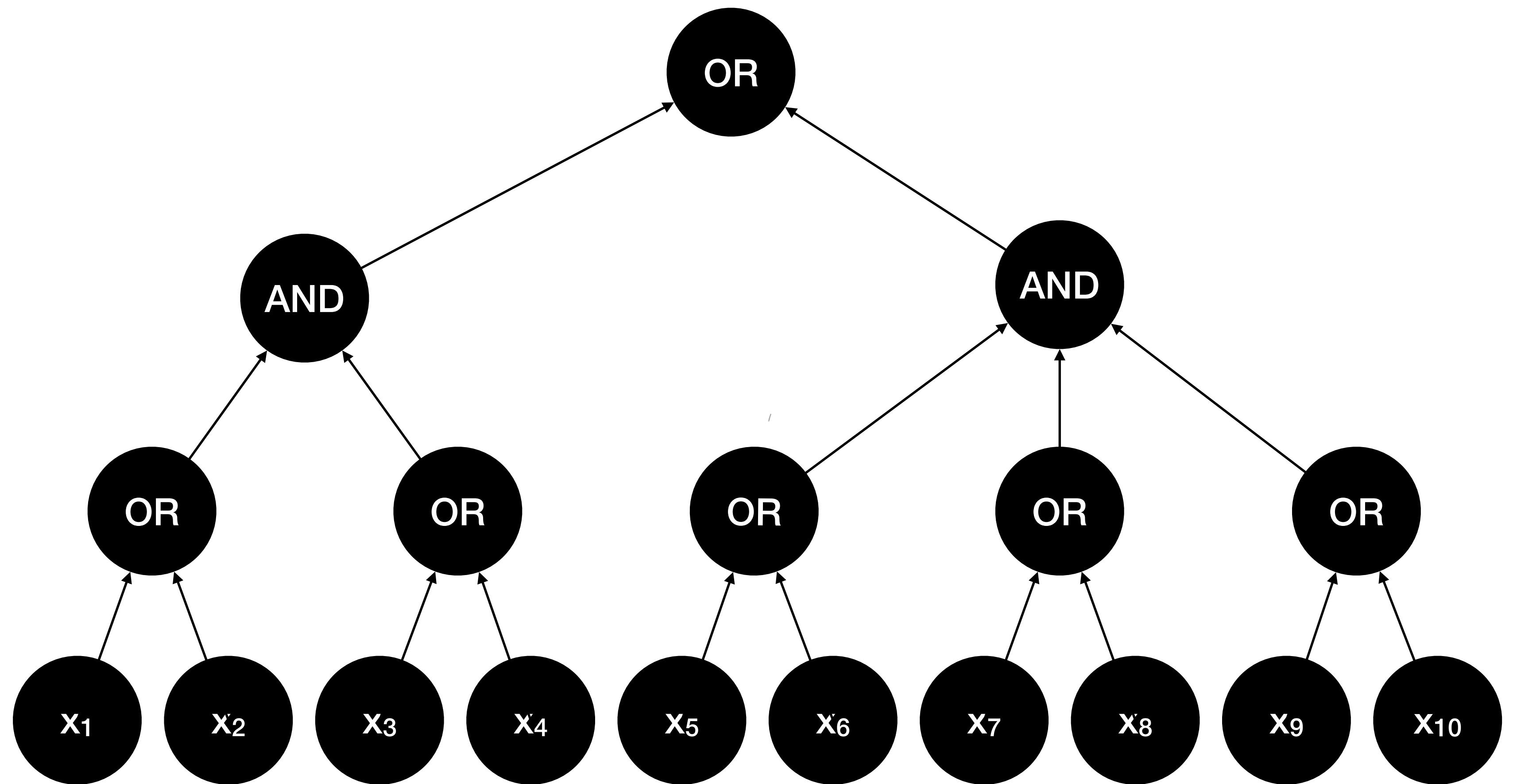


# Depth-3 Circuits for Inner Products

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# $\Sigma_3^2$ : depth-3 bottom fan-in 2 circuits



2-CNFs

# Inner Product

- $\text{IP}_n : \{0,1\}^n \times \{0,1\}^n \rightarrow \{0,1\}$
- $\text{IP}_n(x, y) := \langle x, y \rangle \bmod 2$
- $\text{IP}_9(x, y) = 1$

1	0	0	1	0	0	0	1	0	1
0	0	1	1	0	1	1	0	1	1

$$x \in \{0,1\}^9$$

$$y \in \{0,1\}^9$$

# Why $\Sigma_3^2$ ? Why Inner-Product?

- $\Sigma_3^3$ -complexity of  $\text{IP}_n$  was raised as an open problem by Golovnev, Kulikov, and Williams
- Frankl, Gryaznov, and Talebanfard pointed out the challenge for obtaining  $\Sigma_3^2$ -complexity of  $\text{IP}_n$
- New techniques from related areas

# Our Result

- $\Sigma_3^2$ -complexity of  $\text{IP}_n$  is  $2^{\alpha_n \cdot n}$  where  $0.847\dots \leq \alpha_n \leq 0.965\dots$
- Previously, known:  $\frac{1}{2} \leq \alpha_n \leq 1$

**Upper Bound:**  $\alpha_n \leq 0.965\dots$

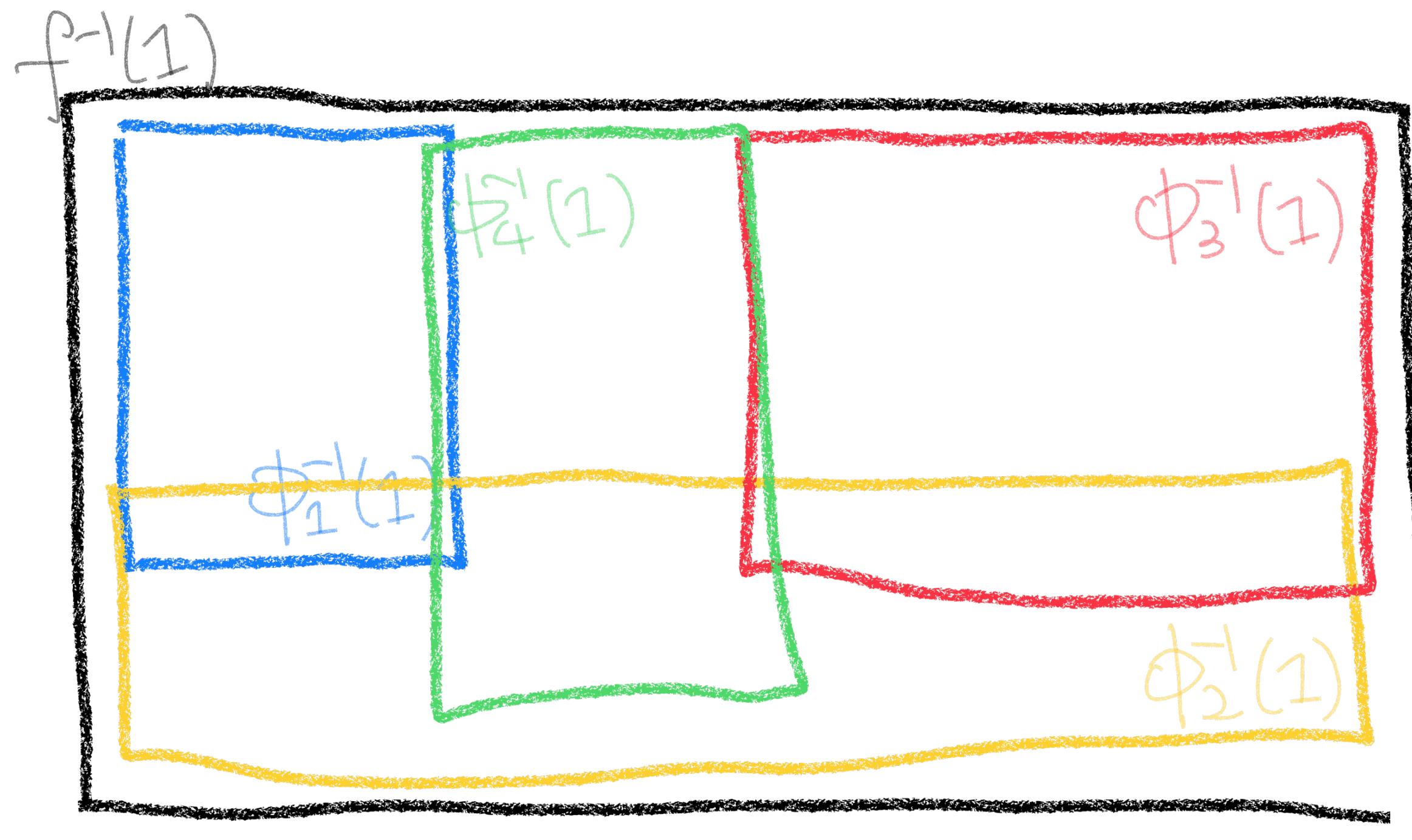
# Upper Bound: $\alpha_n \leq 0.965\dots$

## Fractional Cover [1/4]

- $f : \{0,1\}^n \rightarrow \{0,1\}$
- $C : \Sigma_3^2$ -circuit with top fan-in  $m$  computing  $f$
- $f^{-1}(1) = \bigcup_{i \in [m]} \phi_i^{-1}(1)$
- $\phi_i$ : 2-CNF induced by  $C$

# Upper Bound: $\alpha_n \leq 0.965\dots$

## Fractional Cover [2/4]



# Upper Bound: $\alpha_n \leq 0.965\dots$

## Fractional Cover [3/4]

- $\Phi := \{\phi : 2\text{CNF s.t. } \forall x \in \{0,1\}^n, \phi(x) \leq f(x)\}$
- Goal: Find minimum  $S \subseteq \Phi$  where  $f^{-1}(1) \subseteq \bigcup_{\phi \in S} \phi^{-1}(1)$
- $\Sigma_3^2$ -complexity of  $f \leq O(|S|)$

# Upper Bound: $\alpha_n \leq 0.965\dots$

## Fractional Cover [4/4]

$$\min \sum_{\phi \in \Phi} w_\phi$$

$$\text{subject to } \sum_{\phi \in \Phi} w_\phi \phi(x) \geq 1, \forall x \in f^{-1}(1)$$

$$w_\phi \in [0,1], \forall \phi \in \Phi$$

Note The least top fan-in is captured by the corresponding integer program

Lovasz's fractional cover

$$\Rightarrow \text{Opt} \leq \text{Opt}^* \leq O(n) \cdot \text{Opt}$$

$$w_\phi \in \{0, 1\}$$

# Upper Bound: $\alpha_n \leq 0.965\dots$

## Construct Circuit [1/2]

1	0	0	0	0	0	0	0	0
0	0	1	0	1	1	1	0	1

$x \in \{0,1\}^9$

$y \in \{0,1\}^9$

- $\text{IP}_n(x, y)$  is determined by the parity of “11” patterns

- $p_{11}(x, y) := \frac{1}{n} \left| \{i \in [n] : (x_i, y_i) = (1,1)\} \right|$

- Case on  $p_{11}$  small and  $p_{11}$  large

# Upper Bound: $\alpha_n \leq 0.965\dots$

## Construct Circuit [2/2]

- $p_{11}$  small
  - Hardcode all  $(x, y) \in \text{IP}_n^{-1}(1)$  into a relatively small circuit
- $p_{11}$  large
  - Handle  $(x, y) \in \text{IP}_n^{-1}(0)$  by cancelling out matching pairs: use LP

0	0	1	1	0	0	1	0	0
0	0	1	1	0	0	1	0	0
0	0	1	1	0	0	1	0	0

**Lower Bound:**  $\alpha_n \geq 0.847\dots$

# Lower Bounds: $\alpha_n \geq 0.847\dots$

## Fractional Cover [1/3]

- $\max_{\phi \in \Phi} |\phi^{-1}(1)| \leq M \implies \text{top fan-in of } f \geq \frac{f^{-1}(1)}{M}$

$$f^{-1}(1)$$

$$\begin{aligned} & \phi^{-1}(1) \text{ s.t.} \\ & |\phi^{-1}(1)| = \max_{\phi \in \Phi} |\phi^{-1}(1)| \end{aligned}$$

# Lower Bounds: $\alpha_n \geq 0.847\dots$

## Fractional Cover [2/3]

- Choose a hard distribution  $\mathcal{D}$  over  $f^{-1}(1)$
- Measure  $|\phi^{-1}(1)|$  relative to  $\mathcal{D}$
- $\max_{\phi \in \Phi} \Pr_{x \in \mathcal{D}}[\phi(x) = 1] \leq p \implies \text{top fan-in of } f \geq \frac{1}{p}$

# Lower Bounds: $\alpha_n \geq 0.847\dots$

## Fractional Cover [3/3]

- Hard distribution  $\mathcal{D}_p$  over  $\text{IP}_n^{-1}(1)$ : Binomial distribution where

$$(X, Y) \sim \mathcal{D}_p : \Pr[X_i = 1] = \Pr[Y_i = 1] = p$$

- Max-formulas

$$\forall \phi \in \Phi, \Pr_{(x,y) \sim \mathcal{D}_p}[\phi(x, y) = 1] \leq P_M$$

where  $P_M$  is the acceptance probability of the max-formulas

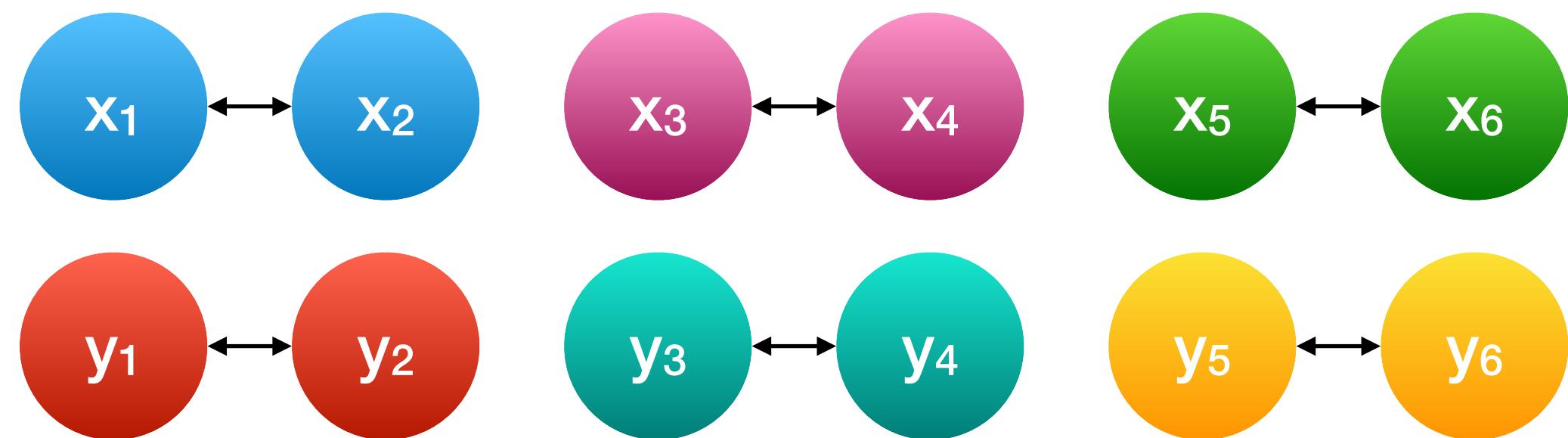
- $\Sigma_3^2$ -complexity of  $\text{IP}_n \geq \frac{1}{P_M}$

# Lower Bounds: $\alpha_n \geq 0.847\dots$

## Max Formulas [1/2]

- $\phi_{\text{Coll}}^{(n)} := \bigwedge_{i \in [n/2]} (x_{2i-1} \leftrightarrow x_{2i}) \wedge (y_{2i-1} \leftrightarrow y_{2i})$

- $\phi_{Nand}^{(n)} := \bigwedge_{i \in [n]} (\neg x_i \vee \neg y_i)$



# Lower Bounds: $\alpha_n \geq 0.847\dots$

## Max Formulas [2/2]

- **Structure lemma:**
  - There exists a finite set  $S_n$  of formulas where
$$\forall \phi \in \Phi, \exists \phi' \in S_n \text{ s.t. } \phi \models \phi' (\phi^{-1}(1) \subseteq (\phi')^{-1}(1))$$
  - There exists  $p^* \in [0,1]$  where
$$\forall \phi \in S_n, \Pr_{\mathcal{D}_{p^*}}[\phi] \leq \max\{\Pr_{\mathcal{D}_{p^*}}[\phi_{Coll}], \Pr_{\mathcal{D}_{p^*}}[\phi_{Nand}]\}$$

**<https://eccc.weizmann.ac.il/report/2023/106/>**